

## Problem 1: Resilient Electric Microgrid

You are installing a microgrid system within a rural area of the power grid. This system lets the rural area to disconnect—or *island*—itself from a nearby city’s power grid during extreme weather events, and power itself via peer-to-peer sharing local renewable energy until the connection can be safely restored. The set of states for the system are:

1. *Interconnected*: The microgrid operates normally, and draws most of its power from the city grid.
2. *Disconnecting*: The microgrid decreases the electric power it draws from the city grid.
3. *Islanded*: The microgrid is using peer-to-peer renewable energy and battery storage sharing.
4. *Connecting*: The microgrid begins reconnecting to the city grid.
5. *Blackout*: There is no power available, and the microgrid cannot operate.

The microgrid behaves in the following way:

1. If an event occurs and a warning is received, the microgrid begins disconnecting from the city grid, until its control system is online, then it moves into *islanded* mode. It stays in *islanded* mode while the event warning is active and it has renewable power available.
2. If an event occurs and the city grid fails to send a warning, the microgrid goes into *blackout*.
3. If the microgrid is in island mode and it runs out of renewable power, it goes into *blackout*.
4. If the microgrid has battery energy storage while in blackout, it returns to *islanded* mode.
5. When the extreme event warning is cleared, the microgrid goes into *connecting* mode.
6. When the control system is offline and the warning is cleared, the microgrid goes into *interconnected* mode.
7. If the microgrid has no battery energy storage, and the warning is still active, it stays in *blackout* until a restoration occurs, directly taking it back to being *interconnected*.

### **Problem 1 (2pts)**

Sketch a state machine diagram for the microgrid system. You don’t need to use any numbers, you can just use words. It might be helpful to define some symbols.

## Problem 2: Return of the Smart Thermostat

Consider the smart thermostat system from Problem Set 1. Equipped with your new knowledge of *sequential logic*, you consider the following inputs for a state machine representation:

1. Input 1: A “too cold” signal  $T_C \in \{0,1\}$ , where  $T_C = 1$  if the house is too cold, that is, temperature  $<$  comfort°.
2. Input 2: A “too hot” signal  $T_H \in \{0,1\}$ , where  $T_H = 1$  if the house is too hot, that is, temperature  $>$  comfort°.

The state machine has two outputs:

1. Heating mode,  $H \in \{0,1\}$ : The HVAC system warms the household
2. Cooling mode,  $C \in \{0,1\}$ : The HVAC system cools the household

Suppose that you want to represent your design via a Moore State Machine. The state machine behaves according to these rules:

1. If the the house is too cold, on the next clock edge, the HVAC should begin heating. If the temperature is no longer too cold, the heater should turn off on the next clock edge.
2. If the the house is too hot, on the next clock edge, the HVAC should begin cooling. If the temperature is no longer too hot, the cooler should turn off on the next clock edge.
3. However, after the heater or cooler turns off, *neither may turn on again* for two clock cycles, even if the temperature is out of the desired range.

**Assume** that the “too hot” and “too cold” signals are never asserted at the same time, that is  $T_C \cdot T_H = 0$  for all time.

### Problem 2 (6pts +2 bonus)

Update the smart thermostat system into a state machine form.

1. Design a Moore state machine that implements the smart thermostat.
2. Suppose the clk signal,  $T_C$ , and  $T_H$  signals evolve over time as shown in Fig. 1. Sketch the heating and cooling signals  $H$  and  $C$  for your smart thermostat system. **You can assume a rising edge trigger.**
3. Propose an improved design of this system that uses the motion detection signal  $M$  from Problem Set 1; specifically, replace the two clock-cycle delay in the problem statement with requiring  $M = 1$ .
4. **[BONUS]** Draw a state-transition table and state machine for your improved design.

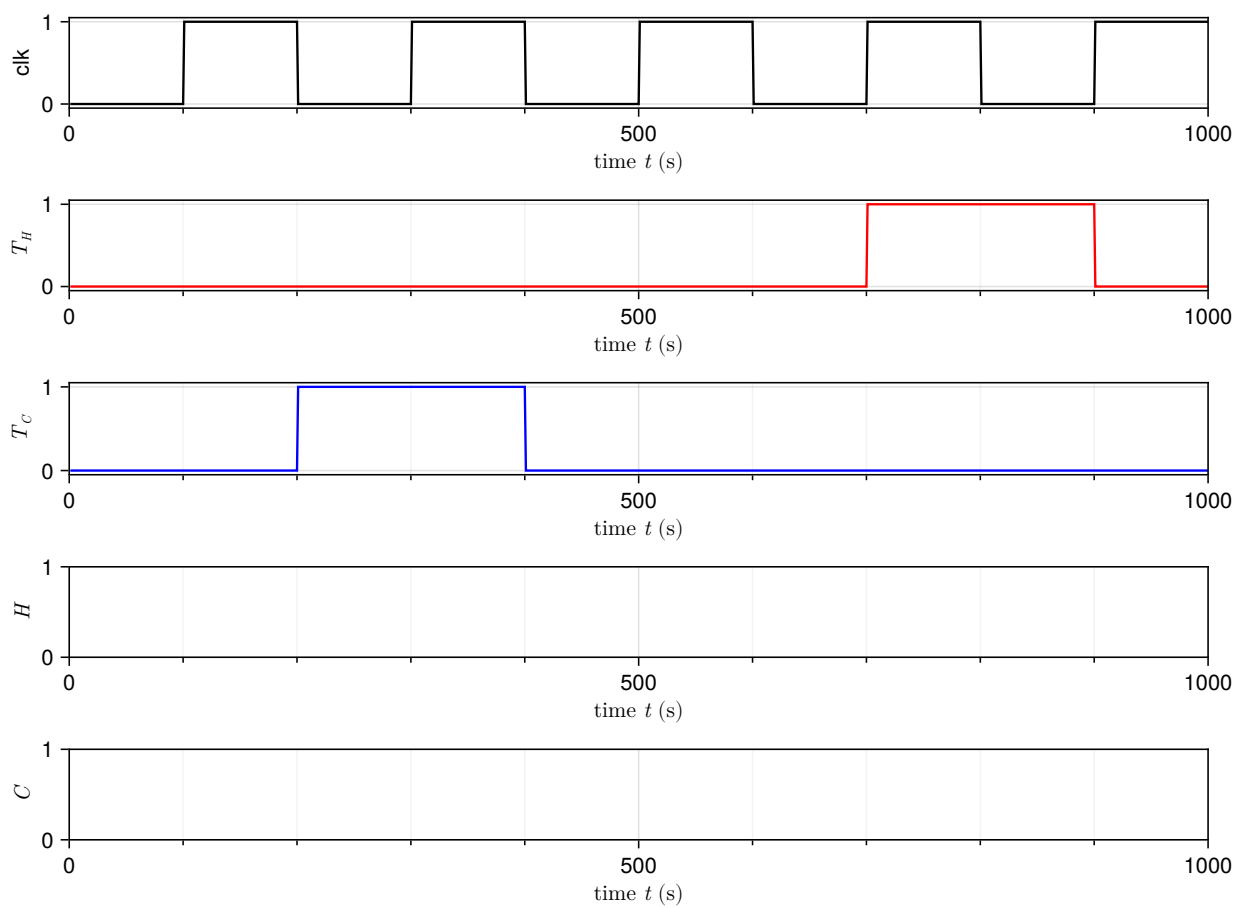


Figure 1: Timing diagram for the smart thermostat

### Problem 3: Mealy and Moore

Consider a Mealy finite state machine with states

$$\mathcal{Q} = \{q : q \in \{0,1\}^2\},$$

and inputs and outputs

$$\mathcal{X} = \{x : x \in \{0,1\}\} \quad \text{and} \quad \mathcal{Y} = \{y : y \in \{0,1\}\},$$

respectively. Suppose that the state machine is described by the following state-transition table.

$q_1$	$q_0$	$x$	$q_1^+$	$q_0^+$	$y$
0	0	0	0	0	0
0	0	1	1	0	1
0	1	0	0	1	0
0	1	1	1	1	1
1	0	0	1	0	1
1	0	1	0	1	0
1	1	0	1	1	1
1	1	1	0	0	0

Table 1: The state-transition table for a Mealy finite state machine with one-dimensional inputs and outputs  $x$  and  $y$ , respectively, and a two-dimensional binary state  $q$ .

#### Problem 3 (6 points)

Explore the connection between Mealy and Moore state machines in the following problems.

1. Draw a Mealy state machine diagram for the state-transition table given in Table 1.
2. Draw a diagram for an equivalent Moore state machine.
3. Draw a state-transition table for your Moore state machine in Part 2.

## Problem 4: Flipping and Flopping

### Problem 4 (2pts)

Consider the sequential logic circuit shown in Fig. 2 using *negative edge-triggered* D flip-flops. Derive logic expressions for  $D_1$ ,  $D_2$ , and  $Z$ , and then sketch their waveforms shown below.

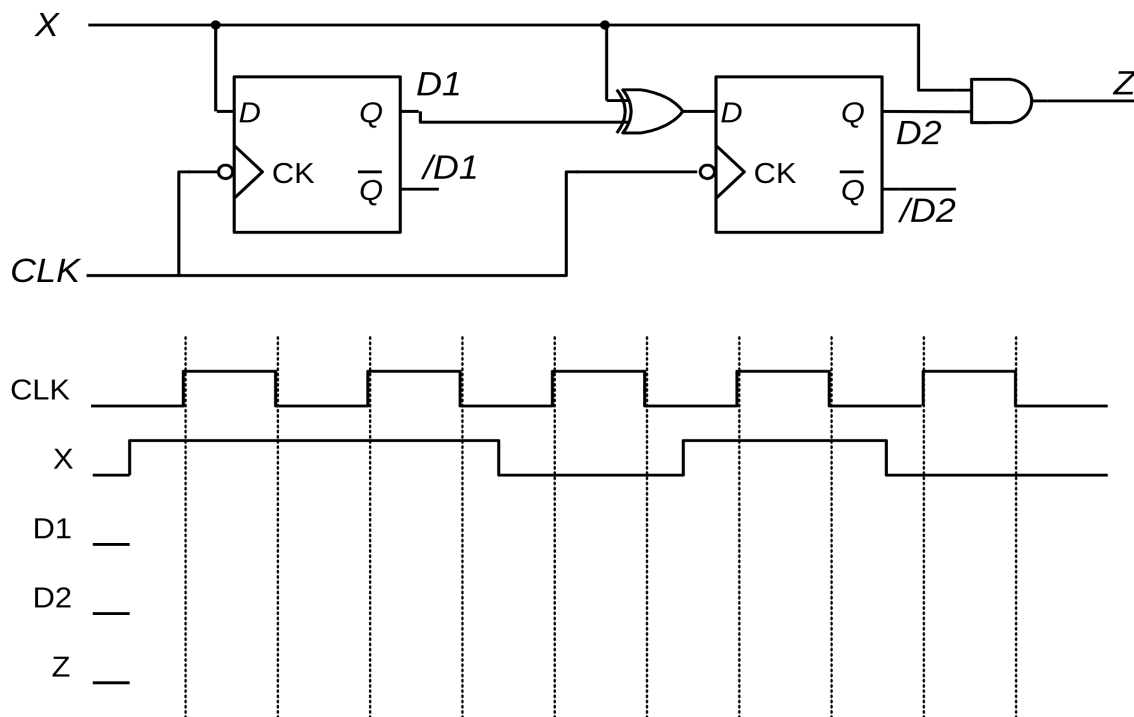


Figure 2: A sequential logic circuit with negative edge-triggered D flip-flops.

### Problem 5: I can't believe it's not a truth table (jk)

current state	next step w/ input 0 $X = 0$	next state w/ input 1 $X = 1$	output
$Q(t)$	$Q(t+1)$	$Q(t+1)$	$Z$
$A$ (initial)	$B$	$D$	0
$B$	$C$	$B$	0
$C$	$B$	$A$	1
$D$	$B$	$C$	0

Table 2: The state-transition table for a positive edge-triggered JK flip flop system.

#### Problem 5 (2pts + 2 bonus)

We can represent the state-transition Table 2 in binary, which we explore in this problem.

1. Suppose you want to design a state machine that implements the state-transition table shown in Table 2 using two positive edge-triggered JK flip-flops. Re-write the state-transition table shown in Table 2 in terms of binary definitions of the states:  $A = (00)_2$ ,  $B = (01)_2$ ,  $C = (10)_2$ , and  $D = (11)_2$ .
2. **[BONUS]:** Derive the state transition function and simplify the expressions for the data signals  $D_1$ ,  $D_0$  associated with the two JK flip flops using K-maps. Optionally sketch your sequential logic circuit design.

## Problem 6: Ups and Downs

Consider the prediction device from Problem 1 of Exam 2. Suppose that you want to upgrade the device to set the tie-breaking signal  $B$  by *counting the number of times* that each predictor  $P_k$  has been right in the past.

This new part of the circuit maintains a *count*  $Q_k$  for each  $k \in \{0, 1, 2, 3\}$ . The circuit has *two inputs* for the stock that are *functions of time*:

1.  $U(t)$ , UP signal: The price of the stock went *up* at the clock cycle  $t$ .
2.  $D(t)$ , DOWN signal: The price of the stock went *down* at the clock cycle  $t$ .

The circuit needs to meet the following requirements for each predictor  $k$ :

1. The circuit does not count, that is,  $Q_k(t) = Q_k(t - 1)$  if  $D(t) = U(t) = 1$  or  $D(t) = U(t) = 0$
2. The circuit should count *up* for  $k$  if predictor  $k$  *was right* at the last clock cycle, that is,  $Q_k(t + 1) = Q_k(t) + 1$ , if the following is true:  $(U(t) = 1 \text{ AND } D(t) = 0 \text{ AND } P_k(t) = 1)$  OR  $(D(t) = 0 \text{ AND } U(t) = 0 \text{ AND } P_k(t) = 0)$ .
3. The circuit should count *down* for  $k$  if predictor  $k$  *was wrong* at the last clock cycle, that is,  $Q_k(t + 1) = Q_k(t) - 1$ , if the following is true:  $(D(t) = 1 \text{ AND } U(t) = 0 \text{ AND } P_k(t) = 1)$  OR  $(U(t) = 1 \text{ AND } D(t) = 0 \text{ AND } P_k(t) = 0)$ .

We will make the following modeling **assumptions**:

1. Assume that the circuit can simply be copied for each predictor  $k$ , so we can design the circuit for a single predictor  $k$ . Therefore, drop the notation subscript  $k$  for the rest of the problem; that is, we will write the signals as  $P_k(t) \stackrel{\text{def}}{=} P(t)$ , and  $Q_k(t) \stackrel{\text{def}}{=} Q(t)$  for convenience.
2. Assume that the circuit can only remember the past 4 trading cycles, so all numbers can be represented as 2-bit binary numbers.

### Problem 6 (6pts + 2 bonus)

Design an electronic counter for tracking the performance of your predictors.

1. **Part 1:** Suppose that  $P(t) = 1$  for all time, (i.e., the predictor is always bullish). Therefore, calculating  $C(t)$  is equivalent to simply summing the UPs and subtracting the DOWNS.
  - (a) Draw a state diagram of this machine. Let the initial state be  $Q(0) = (00)_2$ .
  - (b) Using two positive edge-triggered D flip-flops, implement the circuit, *showing all steps*. You are not required to draw out the circuit, just come up with equations for the flip-flop inputs. *Hint: Use a state transition table, and two K-maps*
  - (c) Sketch the timing diagram shown below in Fig. 3.
2. **Part 2 [BONUS]:** Assume that  $P(t)$  can vary across time. Propose a revision of your design in Part 1 that meets the requirements of the circuit in the problem statement.

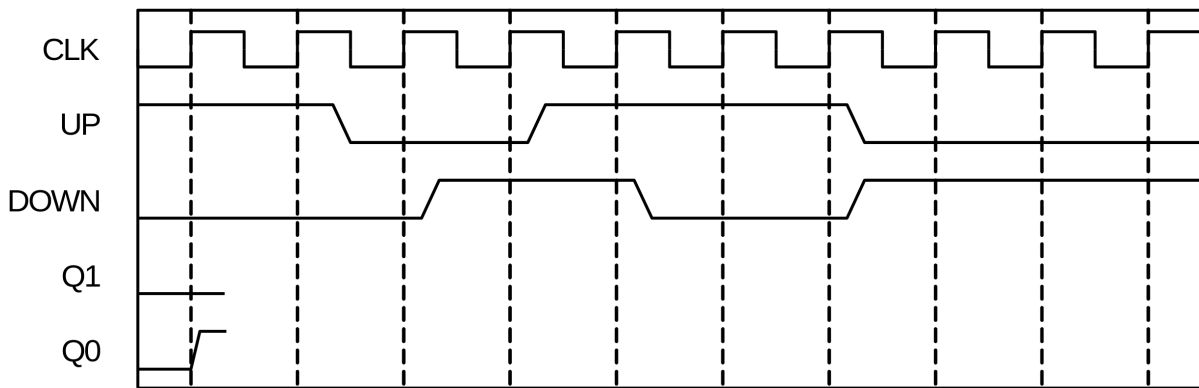


Figure 3: Sketch the prediction accuracy counter